

Enabling Client-Side Crash-Resistance to Overcome Diversification and Information Hiding

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○ Crash-Resistance

Crash-Resistance

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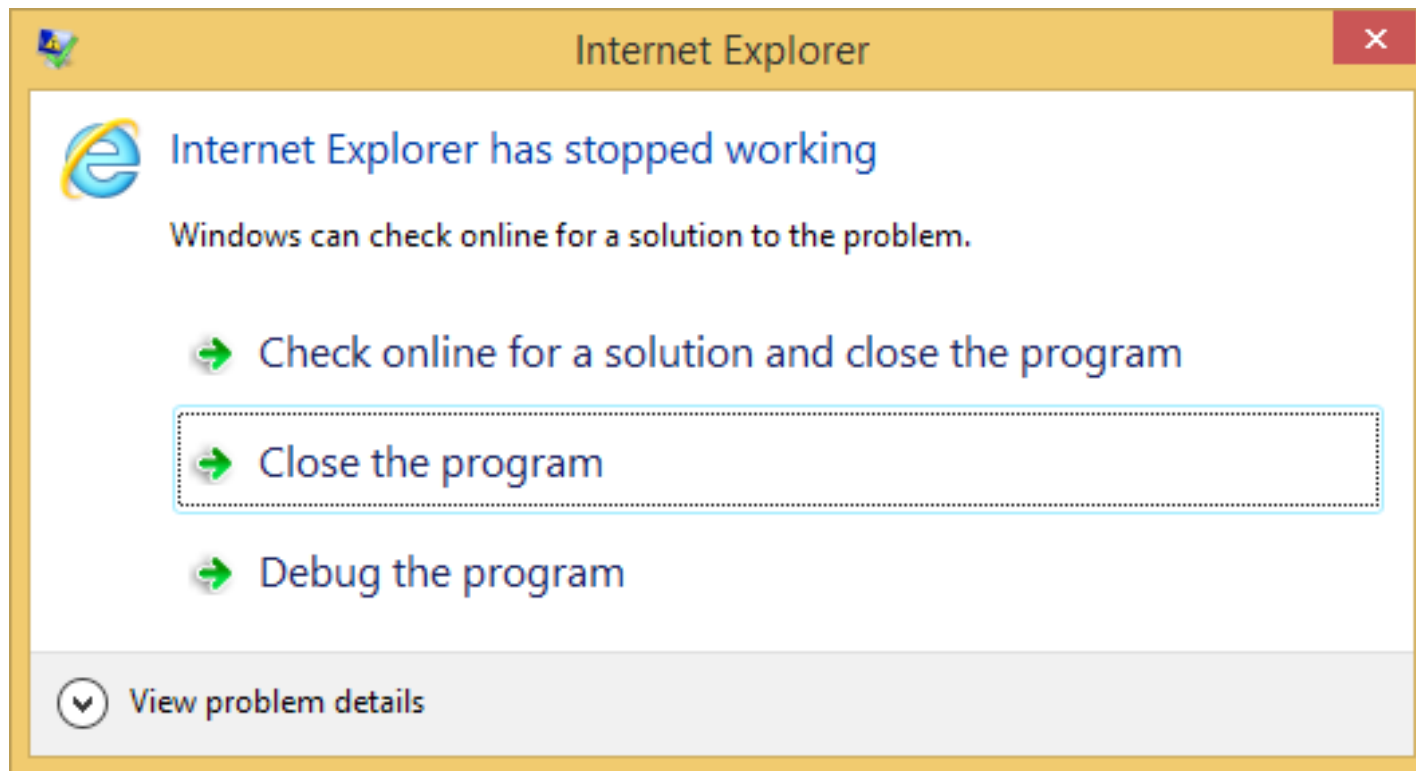
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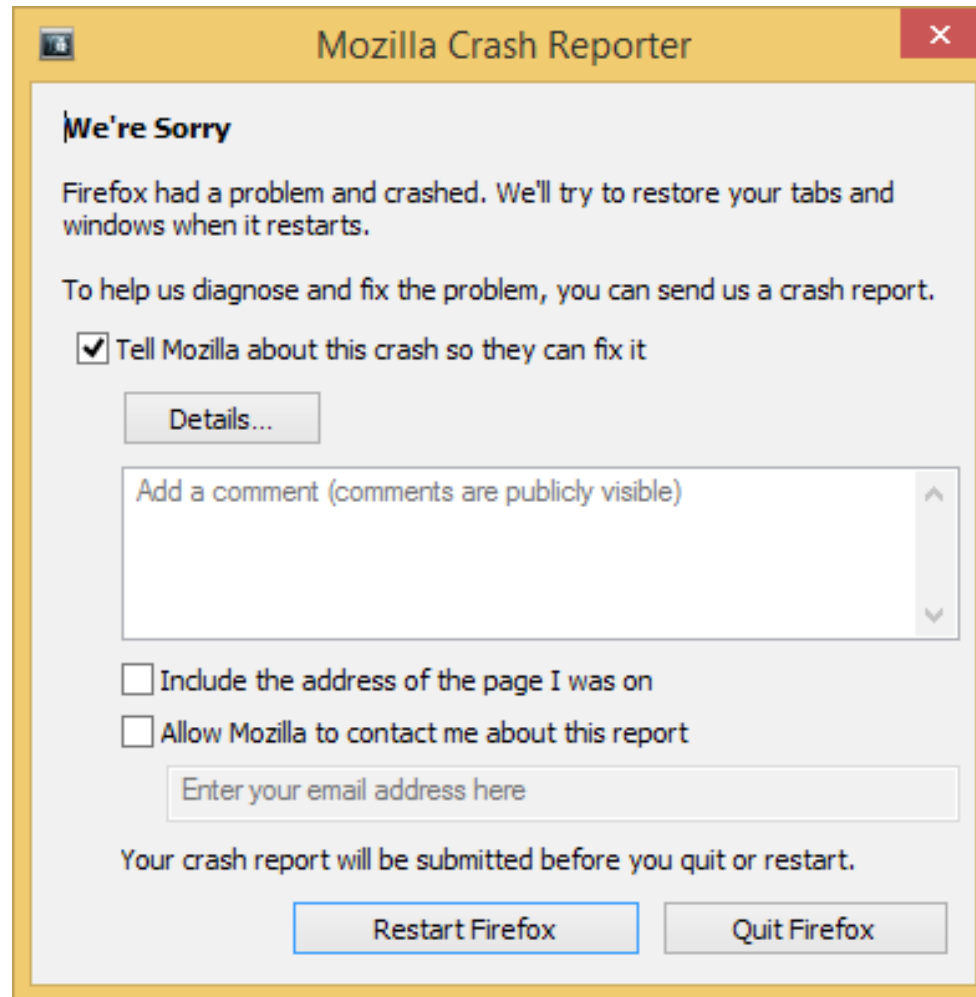
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Instead:
Program runs endlessly

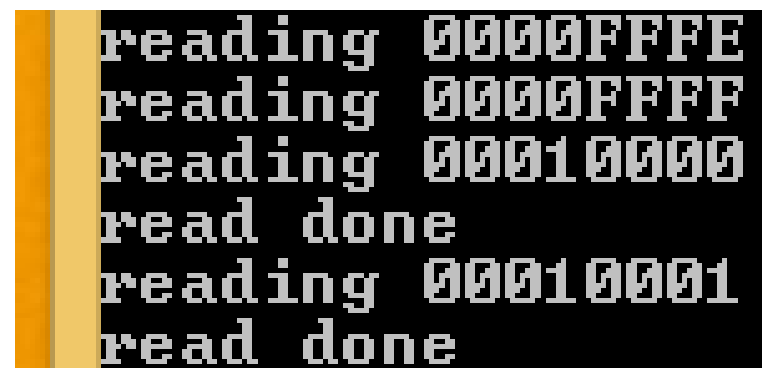
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```
reading 0000FFFE
reading 0000FFFF
reading 00010000
read done
reading 00010001
read done
```

Crash-Resistance

Behind the Scenes

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**Program continues
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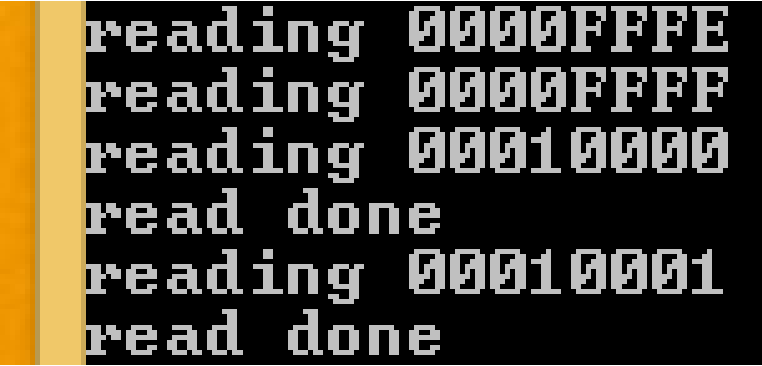
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- It is possible to access memory *more than once* with wrong permissions
 - Client-Side Crash-Resistance is usable as an *attack primitive*

Attacks with Client-Side Crash-Resistance

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 - If *address* is readable, *content* is returned into JavaScript variable
 - On a fault, reset *address* and try reading again

Memory Oracle in Internet Explorer (32-bit)

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Memory Oracle in Mozilla Firefox (64-bit)

- *asm.js* uses exception handling for certain memory accesses
- Modification of *metadata* allows crash-resistant memory queries
- ≈ 700 probes/s (Windows)
- $\approx 18,000$ probes/s (Linux)




Unveiling reference-less hidden memory regions

- memory region is randomized by *ASLR*
- *no references* exist to memory region

First program run



Address space




-  : *readable memory*
-  : *nonreadable memory*
-  : *hidden memory*

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 - Pointer protection metadata: $<$ 1s (Linux 64-bit)

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To mount a control flow hijacking attack, perform *whole function code reuse*

○ Crash-Resistant Oriented Programming (CROP)

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- (4) Read *return data* of system call and proceed to step (1)

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- Defenses that prevent *memory corruption vulnerabilities*, can prevent current crash-resistance primitives

Q & A

References

- [1] Shacham et al. **On the effectiveness of address-space randomization**. *CCS 2004*
- [2] Bittau et al. **Hacking blind**. *Security & Privacy 2014*
- [3] Evans et al. **Missing the Point(er)**. *Security & Privacy 2015*
- [4] <https://www.cve.mitre.org/cgi-bin/cvename.cgi?name=CVE-2015-6161>
- [5] https://bugzilla.mozilla.org/show_bug.cgi?id=1135903