



BGP Origin Authentication



The Problem

- ◆ **Any AS can inject any prefix**
 - ❖ **Mistake (most commonly)**
 - ❖ **Malicious**
- ◆ **Effective DoS attack**
- ◆ **No automated way of excluding bogons**
- ◆ **Need mechanism to differentiate between bogons and legit prefixes**



Different Problems

- ◆ **Anyone can masquerade as another AS**
- ◆ **Anyone can tamper with advertisements**
- ◆ **Valid problems**
- ◆ **Need practical solutions**
- ◆ **Not this talk**



Impractical problems

- ◆ **Compromise of a BGP speaker**
- ◆ **Global Byzantine computations are intractable**
- ◆ **Not this talk**



Authenticate the AS path?

- ◆ **AS path performs two functions**
 - ❖ **Prevents routing & forwarding loops**
 - ❖ **Differentiates between paths**
- ◆ **Attacks**
 - ❖ **DoS**
 - ◆ **Add AS number**
 - ◆ **Delete AS number (causing a loop)**
 - ❖ **Shift traffic**
 - ◆ **Move traffic towards or away**



Authenticate the AS path?

- ◆ **Threat environment**
 - ❖ **Must be a transit ISP**
 - ❖ **Global advertisement provides auditing**
 - ❖ **Transit ISPs can attack the data stream too**
 - ❖ **Hard to hide from traceroute**
- ◆ **Is this a problem worth solving?**
- ◆ **At what cost?**



Our Approach

- ◆ **Encode prefixes in DNS**
- ◆ **Use DNSSEC to provide authentication**
- ◆ **Have BGP look up each prefix in DNS**
- ◆ **Paths to prefixes fall into three classes**
 - ❖ **Authenticated**
 - ❖ **Unauthenticated**
 - ❖ **Authentication failures (bogon)**



The Easy Part: The AS RR

◆ Syntax:

<name> AS <AS number> <prefix length>

◆ Semantics:

- ❖ **The prefix represented in <name> can be advertised with origin <AS number> with the given <prefix length> or longer**



An Example

- ◆ **An AS RR:**

 - 125.128.bgp.in-addr.arpa. AS 47 16**

- ◆ **Prefix 128.125/16 is allocated to AS 47**

- ◆ **Longer prefixes also match!**



On the BGP side

- ◆ **BGP does a lookup for each prefix**
- ◆ **Compare results against each path**
- ◆ **Performance issues:**
 - ❖ **BGP speakers can cache relevant RR's**
 - ❖ **Entire allocation tree fits on secondary storage**
 - ❖ **Cache can persist across reboots**



Fun with BGP

- ◆ **If there's a matching AS RR**
 - ❖ **And the origin doesn't authenticate**
 - ◆ **BOGON!!!**
 - ◆ **Log prefix, origin**
 - ◆ **Select a different path**
 - ◆ **Withdraw it, if it has been advertised**
 - ◆ **Generate SNMP trap, ring bells, send pages, wake the dead, etc.**



More fun with BGP

- ◆ **If there's a matching AS RR**
 - ❖ **And the origin authenticates**
 - ◆ **Authenticated paths may be preferred over unauthenticated paths**
 - ◆ **Authentication has a lifetime**
min TTL of all RRs
 - ◆ **Authentication should be rechecked before lifetime expires**



Even more fun with BGP

- ◆ **If there's no authentication information**
 - ❖ **Paths are unauthenticated**
 - ◆ Paths are useable
 - ◆ Same as today -- eases migration
 - ❖ **Exception: authenticated less-specific prefixes are preferred over unauthenticated more-specific prefixes**



Circular DNS dependency

- ◆ **If there is an authenticated path, it is preferred to an unauthenticated path**
- ◆ **Only the authenticated path is announced**
- ◆ **Transitivity holds: the authenticated path always wins and propagates**
- ◆ **Only holds if domains authenticate the origin**



Migration

- ◆ **Inaction results in the status quo**
- ◆ **Action results in increased protection**
- ◆ **Database configured by address assignors**
- ◆ **Transit providers must deploy new code**
- ◆ **No (intractable) flag days**
- ◆ **Security improves with additional deployment**



Aggregation

- ◆ **How do we deal with aggregates?**
- ◆ **Include aggregates in `bgp.in-addr.arpa`**
- ◆ **Looks just like any other prefix, where the owner is the aggregator**



The Hard Part: DNS

- ◆ How do we encode prefixes and prefix allocation?
- ◆ Awkward on non-octet boundaries
- ◆ Use the classless in-addr hack
- ◆ Root is `bgp.in-addr.arpa.` (or `ipv4.nlri.ietf.org.`, or ... ????)
- ◆ Root is administered by ???



Prefix encoding rules

- ◆ **A name is**

 - <label>.<label> . . . <label>.bgp.in-addr.arpa

- ◆ **Rule 1: Add a label and NS RR for every assignment**

- ◆ **Rule 2: For non-octet assignments:**

 - ❖ **The label is <octet>/<length>**

 - ❖ **Add CNAME records for each octet value in the assignment**



Advantages

- ◆ **Solves 95% of the real problems now**
- ◆ **Tractable amount of computation**
- ◆ **Leverages existing technologies**
- ◆ **Readily implementable**
- ◆ **Scales linearly with the number of paths in the global routing table**
- ◆ **Straightforward migration path**



Forward progress?

- ◆ **We need one global solution**
- ◆ **Debate has not selected an alternative**
- ◆ **Need a practical solution**
- ◆ **Prevent the next incident**
- ◆ **Debate must come to a close soon**
- ◆ **Otherwise:**
 - ❖ **The market will decide**
 - ❖ **After the horse has left the barn**



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